

Solutions - Homework 2

(Due date: October 6th @ 5:30 pm)

Presentation and clarity are very important! Show your procedure!

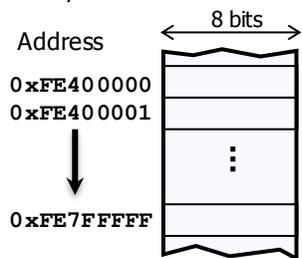
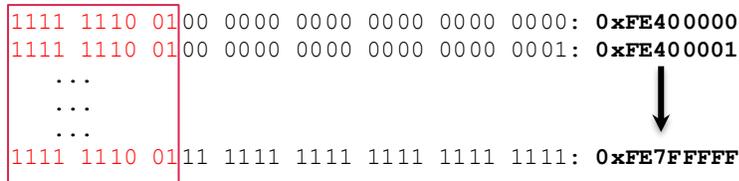
PROBLEM 1 (28 PTS)

- a) What is the minimum number of bits required to represent: (2 pts)
- ✓ 16385 symbols? $\lceil \log_2 16385 \rceil = 15$
 - ✓ Memory addresses from 0 to 131072? $\lceil \log_2(131073) \rceil = 18$

- b) A microprocessor has a 32-bit address line. The size of the memory contents of each address is 8 bits. The memory space is defined as the collection of memory positions the processor can address. (6 pts)
- What is the address range (lowest to highest, in hexadecimal) of the memory space for this microprocessor? What is the size (in bytes, KB, or MB) of the memory space? 1KB = 2^{10} bytes, 1MB = 2^{20} bytes, 1GB = 2^{30} bytes (2 pts)
Address Range: $0x00000000$ to $0xFFFFFFFF$
With 32 bits, we can address 2^{32} bytes, thus we have $2^{2 \cdot 30} = 4GB$ of address space

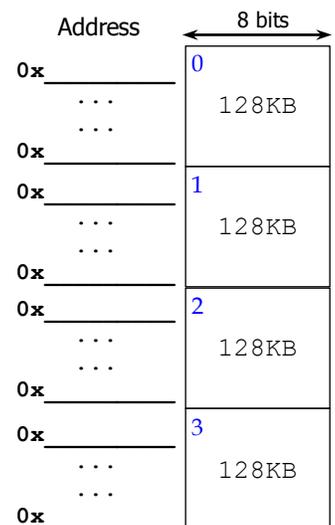
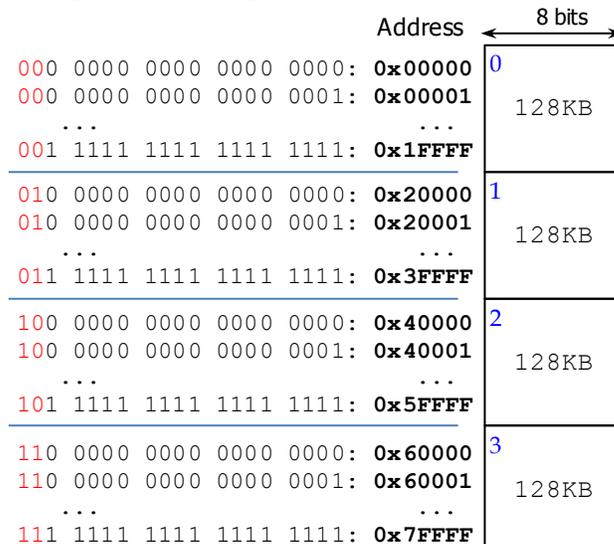
- A memory device is connected to the microprocessor. Based on the size of the memory, the microprocessor has assigned the addresses $0xFE400000$ to $0xFE7FFFFFFF$ to this memory device.
 - What is the size (in bytes, KB, or MB) of this memory device?
 - What is the minimum number of bits required to represent the addresses only for this memory device?

As per the figure, we only need 22 bits for the address in the given range (where the memory device is located). Thus, the size of the memory is $2^{22} = 4 MB$.



- c) A microprocessor has a memory space of 512 KB. Each memory address occupies one byte. (8 pts)
- What is the address bus size (number of bits of the address) of this microprocessor?
Since $512 KB = 2^{19}$ bytes, the address bus size is 19 bits.
 - What is the range (lowest to highest, in hexadecimal) of the memory space for this microprocessor?
With 19 bits, the address range is $0x00000$ to $0x7FFFF$.

- The figure to the right shows four memory chips that are placed in the given positions:
 - Complete the address ranges (lowest to highest, in hexadecimal) for each of the memory chips.



- d) The figure below depicts the entire memory space of a microprocessor. Each memory address occupies one byte. (12 pts)
- What is the size (in bytes, KB, or MB) of the memory space? What is the address bus size of the microprocessor?

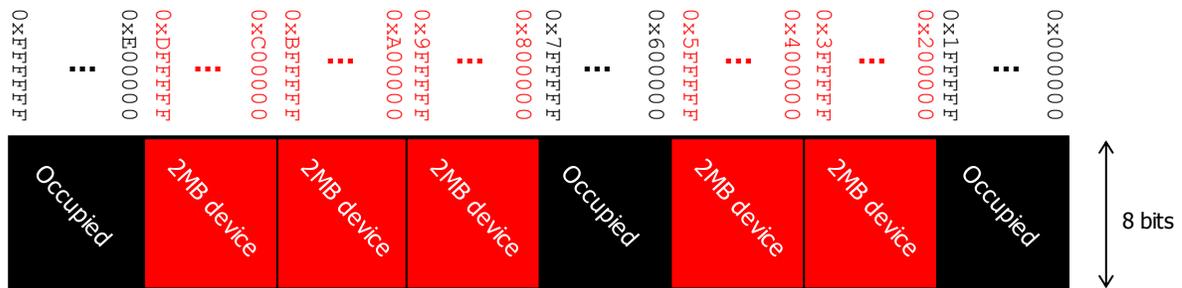
Address space: $0x000000$ to $0xFFFFFFFF$. To represent all these addresses, we require 24 bits. So, the address bus size of the microprocessor is 24 bits. The size of the memory space is then $2^{24} = 16$ MB.

- If we have a memory chip of 2MB, how many bits do we require to address 2MB of memory? (2 pts)

$2\text{ MB} = 2^{21}$ bytes. Thus, we require 21 bits to address only the memory device.

- We want to connect the 2MB memory chip to the microprocessor. Recall that a memory chip must be placed in an address range where every single address shares some MSBs (e.g.: $0x600000$ to $0x7FFFFFFF$). Provide a list of all the possible address ranges that the 2MB memory chip can occupy. You can only use any of the non-occupied portions of the memory space as shown below. (8 pts)

$0x200000$ to $0x3FFFFFFF$
 $0x400000$ to $0x5FFFFFFF$
 $0x800000$ to $0x9FFFFFFF$
 $0xA00000$ to $0xBFFFFFFF$
 $0xC00000$ to $0xDFFFFFFF$



PROBLEM 2 (28 PTS)

- In ALL these problems (a, b, c), you MUST show your conversion procedure. **No procedure = zero points.**
- a) Convert the following decimal numbers to their 2's complement representations: binary and hexadecimal. (8 pts)
 - ✓ -136.6875 , 207.65625 , -128.5078125
 - $207.65625 = 011001111.10101 = 0x0CF.A8$
 - $+128.5078125 = 01000000.1000001 \rightarrow -128.5078125 = 101111111.0111111 = 0xF7F.7E$
- b) Complete the following table. The decimal numbers are unsigned: (6 pts)

Decimal	BCD	Binary	Reflective Gray Code
278	00100111	100010110	110011101
171	000101110001	10101011	11111110
507	010100000111	111111011	100000110
995	100110010101	1111100011	1000010010
217	001000010111	11011001	10110101
84	10000100	01010100	01111110

- c) Complete the following table. Use the fewest number of bits in each case: (14 pts)

Decimal	REPRESENTATION		
	Sign-and-magnitude	1's complement	2's complement
-129	110000001	101111110	101111111
84	01010100	01010100	01010100
-88	11011000	10100111	10101000
0	00	11111	0
-64	11000000	10111111	10000000
-39	1100111	1011000	1011001
-1	11	10	11111

PROBLEM 3 (38 PTS)

a) Perform the following additions and subtractions of the following unsigned integers. Use the fewest number of bits n to represent both operators. Indicate every carry (or borrow) from c_0 to c_n (or b_0 to b_n). For the addition, determine whether there is an overflow. For the subtraction, determine whether we need to keep borrowing from a higher bit. (8 pts)

Example ($n=8$):

✓ $54 + 210$

$$\begin{array}{r}
 \overset{c_8=1}{c_7=1} \quad \overset{c_6=1}{c_5=1} \quad \overset{c_4=0}{c_3=1} \quad \overset{c_2=1}{c_1=0} \quad \overset{c_0=0}{c_0=0} \\
 54 = 0x36 = 00110110 + \\
 210 = 0xD2 = 11010010 \\
 \hline
 \text{Overflow!} \rightarrow 100001000
 \end{array}$$

✓ $77 - 194$

$$\begin{array}{r}
 \text{Borrow out!} \rightarrow \overset{b_8=1}{b_7=0} \quad \overset{b_6=0}{b_5=0} \quad \overset{b_4=0}{b_3=0} \quad \overset{b_2=1}{b_1=0} \quad \overset{b_0=0}{b_0=0} \\
 77 = 0x4D = 01001101 - \\
 194 = 0xC2 = 11000010 \\
 \hline
 00001011
 \end{array}$$

✓ $209 + 68$
✓ $34 + 437$

✓ $129 - 37$
✓ $42 - 98$

$n = 8$ bits

Overflow! $c_8=1$

$$\begin{array}{r}
 \overset{c_8=1}{c_7=1} \quad \overset{c_6=0}{c_5=0} \quad \overset{c_4=0}{c_3=0} \quad \overset{c_2=0}{c_1=0} \quad \overset{c_0=0}{c_0=0} \\
 209 = 0xD1 = 11010001 + \\
 68 = 0x44 = 01000100 \\
 \hline
 277 = 0x115 = 100010101
 \end{array}$$

$n = 8$ bits

No Borrow Out $b_8=0$

$$\begin{array}{r}
 \overset{b_8=0}{b_7=1} \quad \overset{b_6=1}{b_5=1} \quad \overset{b_4=1}{b_3=1} \quad \overset{b_2=0}{b_1=0} \quad \overset{b_0=0}{b_0=0} \\
 129 = 0x81 = 10000001 - \\
 37 = 0x25 = 00100101 \\
 \hline
 92 = 0x5C = 01011100
 \end{array}$$

$n = 9$ bits

No Overflow $c_9=0$

$$\begin{array}{r}
 \overset{c_9=0}{c_8=0} \quad \overset{c_7=0}{c_6=1} \quad \overset{c_5=0}{c_4=0} \quad \overset{c_3=0}{c_2=0} \quad \overset{c_1=0}{c_0=0} \\
 437 = 0x1B5 = 110110101 + \\
 34 = 0x022 = 000100010 \\
 \hline
 471 = 0x1D7 = 111010111
 \end{array}$$

$n = 7$ bits

Borrow out! $b_7=1$

$$\begin{array}{r}
 \overset{b_7=1}{b_6=0} \quad \overset{b_5=0}{b_4=0} \quad \overset{b_3=0}{b_2=0} \quad \overset{b_1=0}{b_0=0} \\
 42 = 0x2A = 0101010 - \\
 98 = 0x62 = 1100010 \\
 \hline
 0x48 = 1001000
 \end{array}$$

b) We need to perform the following operations, where numbers are represented in 2's complement: (24 pts)

✓ $-98 + 256$
✓ $206 + 309$
✓ $-257 + 256$

✓ $105 - 62$
✓ $-127 - 36$
✓ $246 + 31$

For each case:

- ✓ Determine the minimum number of bits required to represent both summands. You might need to sign-extend one of the summands, since for proper summation, both summands must have the same number of bits.
- ✓ Perform the binary addition in 2's complement arithmetic. The result must have the same number of bits as the summands.
- ✓ Determine whether there is overflow by:
 - i. Using c_n, c_{n-1} (carries).
 - ii. Performing the operation in the decimal system and checking whether the result is within the allowed range for n bits, where n is the minimum number of bits for the summands.
- ✓ If we want to avoid overflow, what is the minimum number of bits required to represent both the summands and the result?

$n = 10$ bits

$c_{10} \oplus c_9 = 0$
No Overflow $c_{10}=1$

$$\begin{array}{r}
 \overset{c_{10}=1}{c_9=1} \quad \overset{c_8=0}{c_7=0} \quad \overset{c_6=0}{c_5=0} \quad \overset{c_4=0}{c_3=0} \quad \overset{c_2=0}{c_1=0} \quad \overset{c_0=0}{c_0=0} \\
 -98 = 111001110 + \\
 256 = 010000000 \\
 \hline
 158 = 001001110 \\
 -98+256 = 158 \in [-2^9, 2^9-1] \rightarrow \text{no overflow}
 \end{array}$$

$n = 8$ bits

$c_8 \oplus c_7 = 0$
No Overflow $c_8=1$

$$\begin{array}{r}
 \overset{c_8=1}{c_7=1} \quad \overset{c_6=0}{c_5=0} \quad \overset{c_4=0}{c_3=0} \quad \overset{c_2=0}{c_1=0} \quad \overset{c_0=0}{c_0=0} \\
 105 = 01101001 + \\
 -62 = 11000010 \\
 \hline
 43 = 00101011 \\
 105-62 = 43 \in [-2^7, 2^7-1] \rightarrow \text{no overflow}
 \end{array}$$

n = 10 bits

$C_{10} \oplus C_9 = 1$
Overflow!

$$\begin{array}{r}
 \overset{C_{10}}{0} \quad \overset{C_9}{1} \quad \overset{C_8}{0} \quad \overset{C_7}{1} \quad \overset{C_6}{1} \quad \overset{C_5}{1} \quad \overset{C_4}{1} \quad \overset{C_3}{1} \quad \overset{C_2}{1} \quad \overset{C_1}{0} \quad \overset{C_0}{0} \\
 206 = 0 \ 0 \ 1 \ 1 \ 0 \ 0 \ 1 \ 1 \ 1 \ 0 \ + \\
 309 = 0 \ 1 \ 0 \ 0 \ 1 \ 1 \ 0 \ 1 \ 0 \ 1 \\
 \hline
 1 \ 0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 1 \ 1
 \end{array}$$

206+309 = 515 $\notin [-2^9, 2^9-1]$ → overflow!

To avoid overflow:

n = 11 bits (sign-extension)

$C_{11} \oplus C_{10} = 0$
No Overflow

$$\begin{array}{r}
 \overset{C_{11}}{0} \quad \overset{C_{10}}{0} \quad \overset{C_9}{1} \quad \overset{C_8}{0} \quad \overset{C_7}{1} \quad \overset{C_6}{1} \quad \overset{C_5}{1} \quad \overset{C_4}{1} \quad \overset{C_3}{1} \quad \overset{C_2}{1} \quad \overset{C_1}{0} \quad \overset{C_0}{0} \\
 206 = 0 \ 0 \ 0 \ 1 \ 1 \ 0 \ 0 \ 1 \ 1 \ 1 \ 0 \ + \\
 309 = 0 \ 0 \ 1 \ 0 \ 0 \ 1 \ 1 \ 0 \ 1 \ 0 \ 1 \\
 \hline
 515 = 0 \ 1 \ 0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 1 \ 1
 \end{array}$$

206+309 = 515 $\in [-2^{10}, 2^{10}-1]$ → no overflow

n = 10 bits

$C_{10} \oplus C_9 = 0$
No Overflow

$$\begin{array}{r}
 \overset{C_{10}}{0} \quad \overset{C_9}{0} \quad \overset{C_8}{1} \quad \overset{C_7}{0} \quad \overset{C_6}{0} \quad \overset{C_5}{0} \quad \overset{C_4}{0} \quad \overset{C_3}{0} \quad \overset{C_2}{0} \quad \overset{C_1}{0} \quad \overset{C_0}{0} \\
 -257 = 1 \ 0 \ 1 \ 1 \ 1 \ 1 \ 1 \ 1 \ 1 \ 1 \ + \\
 256 = 0 \ 1 \ 0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 0 \\
 \hline
 -1 = 1 \ 1 \ 1 \ 1 \ 1 \ 1 \ 1 \ 1 \ 1 \ 1
 \end{array}$$

-257+256 = -1 $\in [-2^9, 2^9-1]$ → no overflow

n = 8 bits

$C_8 \oplus C_7 = 1$
Overflow!

$$\begin{array}{r}
 \overset{C_8}{1} \quad \overset{C_7}{0} \quad \overset{C_6}{0} \quad \overset{C_5}{0} \quad \overset{C_4}{0} \quad \overset{C_3}{0} \quad \overset{C_2}{0} \quad \overset{C_1}{0} \quad \overset{C_0}{0} \\
 -127 = 1 \ 0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 1 \ + \\
 -36 = 1 \ 1 \ 0 \ 1 \ 1 \ 1 \ 0 \ 0 \\
 \hline
 0 \ 1 \ 0 \ 1 \ 1 \ 1 \ 0 \ 1
 \end{array}$$

-127-36 = -163 $\notin [-2^7, 2^7-1]$ → overflow!

To avoid overflow:

n = 9 bits (sign-extension)

$C_9 \oplus C_8 = 0$
No Overflow

$$\begin{array}{r}
 \overset{C_9}{1} \quad \overset{C_8}{1} \quad \overset{C_7}{0} \quad \overset{C_6}{0} \quad \overset{C_5}{0} \quad \overset{C_4}{0} \quad \overset{C_3}{0} \quad \overset{C_2}{0} \quad \overset{C_1}{0} \quad \overset{C_0}{0} \\
 -127 = 1 \ 1 \ 0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 1 \ + \\
 -36 = 1 \ 1 \ 1 \ 0 \ 1 \ 1 \ 1 \ 0 \ 0 \\
 \hline
 -163 = 1 \ 0 \ 1 \ 0 \ 1 \ 1 \ 1 \ 0 \ 1
 \end{array}$$

-127-36 = -163 $\in [-2^8, 2^8-1]$ → no overflow

n = 9 bits

$C_9 \oplus C_8 = 1$
Overflow!

$$\begin{array}{r}
 \overset{C_9}{0} \quad \overset{C_8}{1} \quad \overset{C_7}{1} \quad \overset{C_6}{1} \quad \overset{C_5}{1} \quad \overset{C_4}{1} \quad \overset{C_3}{1} \quad \overset{C_2}{0} \quad \overset{C_1}{0} \quad \overset{C_0}{0} \\
 246 = 0 \ 1 \ 1 \ 1 \ 1 \ 0 \ 1 \ 1 \ 0 \ + \\
 31 = 0 \ 0 \ 0 \ 0 \ 1 \ 1 \ 1 \ 1 \ 1 \\
 \hline
 277 = 1 \ 0 \ 0 \ 0 \ 1 \ 0 \ 1 \ 0 \ 1
 \end{array}$$

246+31 = 277 $\notin [-2^8, 2^8-1]$ → overflow!

To avoid overflow:

n = 10 bits (sign-extension)

$C_{10} \oplus C_9 = 0$
No Overflow

$$\begin{array}{r}
 \overset{C_{10}}{0} \quad \overset{C_9}{0} \quad \overset{C_8}{1} \quad \overset{C_7}{1} \quad \overset{C_6}{1} \quad \overset{C_5}{1} \quad \overset{C_4}{1} \quad \overset{C_3}{1} \quad \overset{C_2}{0} \quad \overset{C_1}{0} \quad \overset{C_0}{0} \\
 246 = 0 \ 0 \ 1 \ 1 \ 1 \ 1 \ 0 \ 1 \ 1 \ 0 \ + \\
 31 = 0 \ 0 \ 0 \ 0 \ 0 \ 1 \ 1 \ 1 \ 1 \ 1 \\
 \hline
 277 = 0 \ 1 \ 0 \ 0 \ 0 \ 1 \ 0 \ 1 \ 0 \ 1
 \end{array}$$

246+31 = 277 $\in [-2^9, 2^9-1]$ → no overflow

- c) Get the multiplication results of the following numbers that are represented in 2's complement arithmetic with 4 bits. (6 pts)
✓ 1001x1011, 1010x0101, 1110x0110.

$$\begin{array}{l}
 \begin{array}{r}
 1 \ 0 \ 0 \ 1 \times \\
 1 \ 0 \ 1 \ 1 \\
 \hline
 0 \ 1 \ 1 \ 1 \\
 0 \ 0 \ 0 \ 0 \\
 0 \ 1 \ 1 \ 1 \\
 \hline
 0 \ 1 \ 0 \ 0 \ 0 \ 1 \ 1 \\
 \downarrow \\
 0 \ 1 \ 0 \ 0 \ 0 \ 1 \ 1
 \end{array}
 \quad
 \begin{array}{r}
 1 \ 0 \ 1 \ 0 \times \\
 0 \ 1 \ 0 \ 1 \\
 \hline
 0 \ 1 \ 1 \ 0 \\
 0 \ 0 \ 0 \ 0 \\
 0 \ 1 \ 1 \ 0 \\
 \hline
 0 \ 1 \ 1 \ 1 \ 1 \ 0 \\
 \downarrow \\
 1 \ 0 \ 0 \ 0 \ 1 \ 0
 \end{array}
 \quad
 \begin{array}{r}
 1 \ 1 \ 1 \ 0 \times \\
 0 \ 1 \ 1 \ 0 \\
 \hline
 0 \ 0 \ 1 \ 0 \\
 0 \ 0 \ 0 \ 0 \\
 0 \ 0 \ 1 \ 0 \\
 0 \ 0 \ 1 \ 0 \\
 \hline
 0 \ 0 \ 1 \ 1 \ 0 \ 0 \\
 \downarrow \\
 1 \ 1 \ 0 \ 1 \ 0 \ 0
 \end{array}
 \end{array}$$

PROBLEM 4 (6 PTS)

- Complete the timing diagram (signals *DO* and *DATA*) of the following circuit. The circuit in the blue box is a 4-bit Binary to Gray Decoder. For example, if $T=1100$, then $DO=1010$.

